



The simplest model

Regeneration point  
approach

More complex  
equipment  
replacement problem

# Lecture 7

## Equipment Replacement Problem

*MATH3220 Operations Research and Logistics*  
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# Agenda

- 1 The simplest model
- 2 Regeneration point approach
- 3 More complex equipment replacement problem



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## The Simplest Model

### - Problem

Our basic problem concerns a type of machine (perhaps an auto-mobile) which deteriorates with age, and make decisions about when to replace the incumbent machine, when to replace its replacement, etc., so as to minimize the total cost during the next  $N$  years.

### - Assumption

- We must own such a machine during each of the  $N$  time periods (say years).
- $y$  is the age of the machine when we start the process.
- $c(i)$  is the cost of operating for one year a machine which is of age  $i$  at the start of the year.
- $p$  is the price of a new machine (of age 0).
- $t(i)$  is the trade-in value received when a machine which is of age  $i$  at the start of a year is traded for a new machine at the start of the year.
- $s(i)$  is the salvage value received for a machine that has just turned age  $i$  at the end of year  $N$ .



## Dynamic Programming Model

- (i) OPTIMAL VALUE FUNCTION:  $S(x, k)$  is the minimum cost of owning a machine from year  $k$  through  $N$ , starting year  $k$  with a machine just turned age  $x$ , for  $k = 1, 2, \dots, N$ ;  $x = 1, 2, \dots, k - 1, y + k - 1$  when  $k > 1$ ; and  $x = y$  when  $k = 1$ . Here  $y$  is the age of the starting machine.
- (ii) RECURRENCE RELATION:
- $$S(x, k) = \text{Min} \begin{cases} \text{buy} : p - t(x) + c(0) + S(1, k + 1), \\ \text{keep} : c(x) + S(x + 1, k + 1). \end{cases}$$
- (iii) OPTIMAL POLICY FUNCTION:  $P(x, k) = B$  (buy) if buy is cheaper than keep in the recurrence relation, and  $P(x, k) = K$  (keep) if otherwise.
- (iv) BOUNDARY CONDITION:  $S(x, N + 1) = -s(x)$  for  $x = 1, 2, \dots, N$  and  $y + N$ .
- (v) ANSWER SOUGHT:  $S(y, 1) =$  the minimum cost.



## Example

As an example, consider the following equipment replacement problem:

$$N = 5$$

$y$ (the age of the incumbent machine at the start of year 1) = 2

$$c(0) = 10, c(1) = 13, c(2) = 20, c(3) = 40, c(4) = 70,$$

$$c(5) = 100, c(6) = 100;$$

$$p = 50;$$

$$t(1) = 32, t(2) = 21, t(3) = 11, t(4) = 5, t(5) = 0, t(6) = 0;$$

$$s(1) = 25, s(2) = 17, s(3) = 8, s(4) = 0, s(5) = 0, s(7) = 0.$$

Note that we do not need  $s(6)$ , as there is no chance that the car will be of six years old at the end of fifth year.



## Example

The DP computations are summarized in the following table.

$k$		$x$						
		1	2	3	4	5	6	7
6	$S(x, k)$	-25	-17	-8	0	0	-	0
5	keep	-4	12	40	70	-	100	
	buy	3	14	24	30	-	35	
	$S(x, k)$	-4K	12K	24K	30B	-	35B	
4	keep	25	44	70	-	135		
	buy	24	35	45	-	56		
	$S(x, k)$	24B	35B	45B	-	56B		
3	keep	48	65	-	126			
	buy	52	63	-	79			
	$S(x, k)$	48K	63B	-	79B			
2	keep	76	-	119				
	buy	76	-	97				
	$S(x, k)$	76KB	-	97B				
1	keep	-	117					
	buy	-	115					
	$S(x, k)$	-	115B					

We see that the minimum cost is 115; and the optimal sequence of decisions is *BKBBK* or *BBKBBK*, where *B* is buy and *K* is keep.



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2	keep	76	-	119				
	buy	76	-	97				
	$S(x, k)$	<u>76KB</u>	-	97B				
1	keep	-	117					
	buy	-	115					
	$S(x, k)$	-	<u>115B</u>					

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## Exercise

How many additions and how many comparisons, as a function of the duration of the process  $N$ , are required?



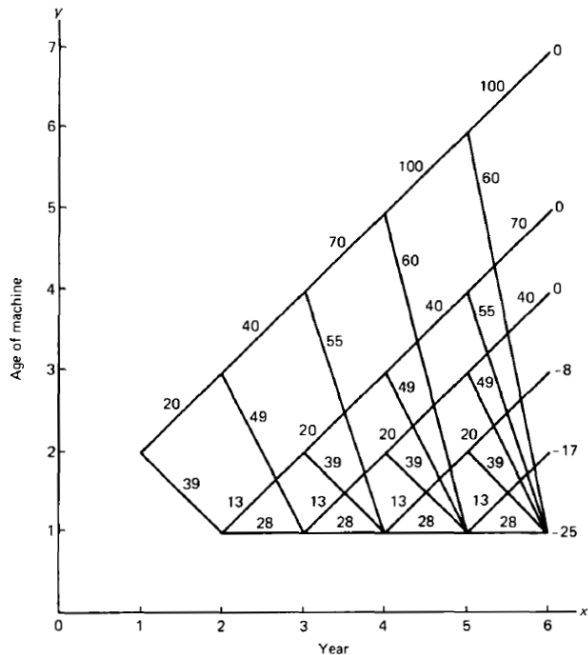
## Shortest-path Representation of the Problem

Almost all dynamic-programming problems can be thought of as problems seeking the minimum-cost path (generally in more than two dimensions and therefore generally not easily drawn).

Letting the  $x$  axis denote the year and the  $y$  axis represent the age of the machine, we start at  $(1, y)$ . The “buy” decision takes us to  $(2, 1)$  at an arc cost of  $p - t(y) + c(0)$  and “keep” leads us to  $(2, y + 1)$  with an arc cost of  $c(y)$ . The same reasoning applies at each point.



# Shortest-path Representation of the Problem

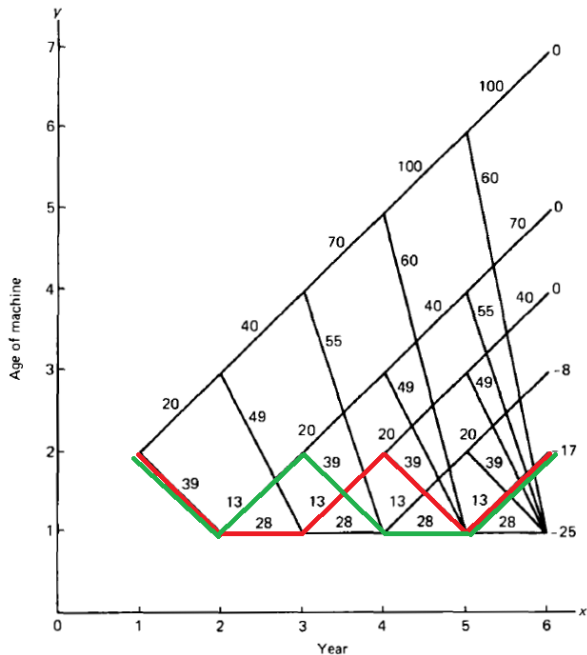


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# Shortest-path Representation of the Problem



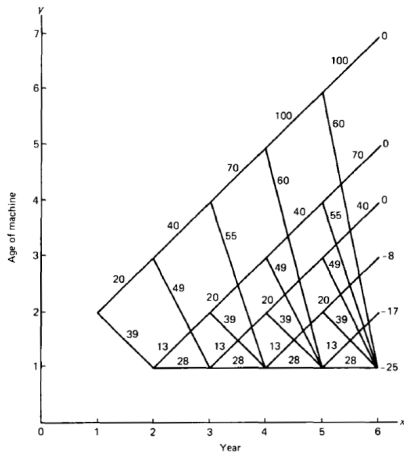
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## Regeneration Point Approach

Unlike many path problem, for the equipment replacement problem, we can be sure in advance that all paths, but one, eventually return at least once to a vertex on the horizontal line  $y = 1$ . When we return to  $y = 1$  the process is said to "regenerate" itself and we can ask the question, "Given an initial situation, when shall we make our first purchase?"





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(i) OPTIMAL VALUE FUNCTION:  $S(i)$  = the minimum attainable cost for the remaining process given we start year  $i$  with a one-year-old machine.

(ii) RECURRENCE RELATION:

$$S(i) = \min \left[ \begin{array}{l} \sum_{k=1}^{N-(i-1)} c(k) - s(N-i+2) \\ \min_{j=i, \dots, N} \left\{ \sum_{k=1}^{j-i} c(k) + p - t(j-i+1) + c(0) + S(j+1) \right\} \end{array} \right]$$

(iii) OPTIMAL POLICY FUNCTION:

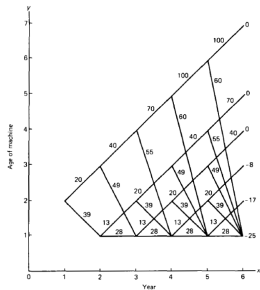
$P(x, k) = \textit{Keep until end}$  if the value in the first row in the recurrence relation is less than the value in the second row, and

$P(x, k) = \textit{Buy at the start of year } j'$  if the minimum value is obtained in the second row in the recurrence relation at  $j = j'$ .

(iv) BOUNDARY CONDITION:  $S(N+1) = -s(1)$

(v) ANSWER SOUGHT:  $S(1)$  ???

Using the procedure to solve for the data shown in the figure, we get



$$S(6) = -25;$$

$$S(5) = \min \left[ \frac{13 - 17}{28 + S(6)} \right] = -4,$$

$P(5) =$  keep until end;

$$S(4) = \min \left[ \frac{13 + 20 - 8}{28 + S(5)}, 13 + 39 + S(6) \right] = 24,$$

$P(4) =$  buy at the start of year 4;

$$S(3) = \min \left[ \frac{13 + 20 + 40}{28 + S(4)}, \frac{13 + 39 + S(5)}{13 + 20 + 49 + S(6)} \right] = 48,$$

$P(3) =$  buy at the start of year 4;

$$S(2) = \min \left[ \frac{13 + 20 + 40 + 70}{28 + S(3)}, \frac{13 + 39 + S(4)}{13 + 20 + 49 + S(5)}, \frac{13 + 20 + 40 + 55 + S(6)}{13 + 20 + 40 + 55 + S(6)} \right] = 48,$$

$P(2) =$  buy at the start of either year 2 or 3.



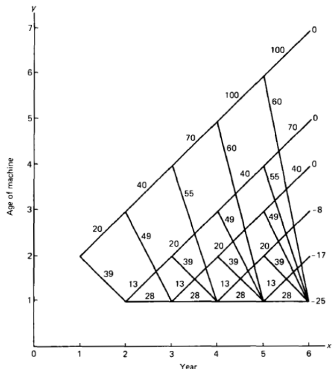


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= 115



The answer is:

$$\min \left[ \begin{array}{l} 20 + 40 + 70 + 100 + 100 \\ 39 + S(2), \\ \hline 20 + 49 + S(3), \\ 20 + 40 + 55 + S(4), \\ 20 + 40 + 70 + 60 + S(5), \\ 20 + 40 + 70 + 100 + 60 + S(S) \end{array} \right]$$

$P(1)$  = buy at start of year 1



## More Complex Equipment Replacement Problem

- In the above equipment-replacement problem, one additional decision is available, namely, “*overhaul*”.
- An overhauled machine is better than one not overhauled, but not as good as a new one.
- Let us further assume that performance depends on the *actual age of equipment* and on the *number of years since last overhaul*, but is independent of when and how often the machine was overhauled prior to its last overhaul.



The known data are

$k$  = the  $k$ th year;

$i$  = a machine's current age;

$j$  = age at last overhaul;

$e(k, i, j)$  = cost of exchanging a machine of age  $i$ , last overhauled at age  $j$  for a new machine at the start of year  $k$ ;

$c(k, i, j)$  = operating cost during year  $k$  of a machine of age  $i$  and last overhauled at age  $j$ ;

$o(k, i)$  = cost of overhauling a machine of age  $i$  at the beginning of year  $k$ ;

$s(i, j)$  = salvage value at the end of year  $N$  of a machine which has just become age  $i$  and last overhauled at  $j$  years ago.

If  $j = 0$ , then the machine has never been overhauled.





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Using the consultant's approach, we can see that, if we want to pursue an optimal replacement policy, the minimal information needed at the start of the  $k$ th year is the age of the car and how long ago it has been overhauled. Thus we have the following optimal value function:

$f(k, i, j)$  = the minimum cost during the remaining years given we start year  $k$  with a machine of age  $i$  and last overhauled at age  $j$ .

The recurrence relation is:

$$f(k, i, j) = \min \left[ \begin{array}{l} \text{Replace: } e(k, i, j) + c(k, 0, 0) + f(k + 1, 1, 0) \\ \text{Keep: } \quad c(k, i, j) + f(k + 1, i + 1, j) \\ \text{Overhaul: } o(k, i) + c(k, i, i) + f(k + 1, i + 1, i) \end{array} \right]$$

and the boundary condition is

$$f(N + 1, i, j) = -s(i, j).$$



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For  $k = N$ , assuming the incumbent machine is new, we must compute  $f$  for  $i = 1, 2, \dots, N - 1$  and  $j = 0, 1, 2, \dots, i - 1$ .

This involves  $N - 1$  evaluations of three decisions for  $i = N - 1, N - 2$  for  $i = N - 2, \dots$ , and 1 for  $i = 1$ .

$\Rightarrow$  A total of  $(N - 1)N/2$  such evaluations.

For  $k = N - 1$ , we have  $(N - 2)(N - 1)/2$  such evaluations.

For  $k = N - 2$ , we have  $(N - 3)(N - 2)/2$  such evaluations.

$\vdots$

$\Rightarrow$  The total number is precisely

$$\sum_{i=2}^N (i - 1)i/2 + 1$$

or, approximately

$$\sum_{i=1}^N i^2/2 \approx N^3/6$$

Consequently, the total number of operations is roughly  $N^3$  since each evaluation of the right-hand side of the recurrence relation requires a total of seven additions and comparisons.